Optimized Hardware Configuration for High Performance Computing Systems

Scott Hutchison Department of Computer Science Kansas State University Manhattan, KS 66505, USA email: scotthutch@ksu.edu

Mitchell Neilsen Department of Computer Science Kansas State University Manhattan, KS 66505, USA email: neilsen@ksu.edu Daniel Andresen Department of Computer Science Kansas State University Manhattan, KS 66505, USA email: dan@ksu.edu William Hsu Department of Computer Science Kansas State University Manhattan, KS 66505, USA email: bhsu@ksu.edu

Benjamin Parsons High Performance Computing Modernization Program Engineering Research and Development Center Vicksburg, MS 39180, USA email: ben.s.parsons@erdc.dren.mil

Abstract—When faced with upgrading or replacing High Performance Computing or High Throughput Computing systems, system administrators can be overwhelmed by hardware options. Servers come with various configurations of memory, processors, and hardware accelerators, like graphics cards. Differing server capabilities greatly affect their performance and their resulting cost. For a fixed budget, it is often difficult to determine what server package composition will maximize the performance of these systems once they are purchased and installed. This research uses simulation to evaluate the performance of different server packages on a set of jobs, and then trains a machine learning model to predict the performance of un-simulated server package compositions. In addition to being orders of magnitude faster than conducting simulations, this model is used to power a recommender system that provides a precision@50 of 92%.

Index Terms—HPC; Procurement Optimization; Recommender system; XGBoost.

I. INTRODUCTION

When faced with upgrading or expanding a High Performance Computing (HPC) or High Throughput Computing (HTC) system, administrators of these systems can be overwhelmed by options. It is a challenging task to get the best performance for a fixed budget. Server capabilities (i.e., number and types of processors, amount of memory, and number and types of Graphics Processing Units (GPU) or other hardware accelerators) greatly affect their costs, and for a fixed spending ceiling, it is desirable to get the "best bang for your buck." For an HPC system, an optimal server package composition is dictated by its typical use. For instance, if many users rely upon a GPU-accelerated application or library, a higher GPU count may be desirable, even if this means fewer servers can be purchased. With many factors to consider, HPC administrators often rely upon their preferences, intuition, and experience to inform procurement decisions. This research uses historical job data from an HPC system, a discrete event simulator, and a machine learning model to power a recommender system, which can help inform a hardware procurement decision. These techniques provide additional information to HPC system administrators about which set of budget-constrained hardware minimizes wait time for users' jobs, and provides quantifiable support for procurement decisions when upgrading or expanding existing HPC infrastructure. The contributions of this work can be summarized as follows:

- 1) A data set consisting of roughly 12,700 HPC scheduling simulations, each with a different HPC server set
- 2) An optimized XGBoost regression model for predicting average wait time when given a composition of servers
- A recommender system with precision@50=92%, which can inform hardware procurement decisions

This paper is laid out as follows: Section II provides additional background on the problem and describes similar work done by others, Section III provides the methodology and some implementation details, Section IV provides details of formulas for metric calculations, Section V provides the results of the experiments, and Sections VI provides our final conclusions.

II. BACKGROUND AND RELATED WORKS

The Open Science Grid (OSG) [1] [2] is a worldwide collaboration that offers distributed computing for scientific research. In the central United States, one of the organizations contributing resources to the OSG is the Great Plains Augmented Regional Gateway to the Open Science Grid (GP-ARGO) [3]. In part, GP-ARGO receives funding through governmental grants. These grants are often used to procure new equipment to expand or improve the capabilities of GP-ARGO's participating organizations. Consequentially, there is a fixed budget ceiling for HPC equipment procurement, and the administrator's goal is to purchase new equipment that will maximize computational performance for our typical applications while ensuring costs remain under the fixed grant budget. The research question for this work is as follows: for a

planned HPC expansion, can experimental simulation provide an optimal set of hardware under a given budget that will minimize job wait time?

The challenge of optimal hardware procurement is not exclusive to our organization. Similar work was done by Evans et al. [4]. They collected benchmarks for various software applications on different hardware to optimize the ratio of Central Processing Unit (CPU) and GPU architectures for HPC jobs. Their work is similar to ours, but we took a different approach by using a scheduling simulator to evaluate the performance of a set of jobs that were actually submitted to an HPC system. We are solving a very similar problem as Evans et al., but using a different approach to arrive at an optimal hardware configuration.

Other researchers have attempted to optimize for a particular application, such as the work Kutzner et al. [5] did to improve the utilization of GPU nodes when using GROMACs. Although these techniques are not without their merits for HPC systems that run a large number of homogeneous applications, users of the GP-ARGO HPC systems run a wide variety of jobs and applications. A more broad scheduler-based optimization was more appropriate for our application.

Various public HPC workloads exist [6], and have been used by HPC researchers in the past. However, as we are attempting to identify and evaluate new hardware for a specific HPC system, log data from that HPC system was utilized as the workload for this research.

Different scheduling applications like SLURM, HTCondor, or PBS, operate on HPC systems and perform the function of assigning HPC resources to jobs. This job to machine assignment task is as an extension of the online bin packing problem [7]. For the bin packing problem, the goal is to pack a sequence of items with sizes between 0 and 1 into as few bins of size 1 as possible. Each job specifies the resources requested (the object sizes), and each HPC machine has a certain amount of available resources (the bins with their respective sizes). The scheduler is given the task to meet job requirements by assigning them to HPC nodes (pack the objects into the available bins) as efficiently as possible. This is an online problem as new jobs are submitted over time to the scheduler. The best fit bin packing (BFBP) algorithm has been shown by Dosa and Sgall [8] to use at most |1.7OPT| bins, ensuring this algorithm will provide a reasonably close to optimal average wait time when it is used as an HPC job scheduling algorithm. Since scheduling algorithms vary between applications, most being highly customizable, and others being proprietary, a discrete event simulator utilizing the BFBP algorithm served as a stand-in for our scheduling application in an attempt to make it more universally applicable. The BFBP scheduling algorithm is described in Figure 1.

Although various HPC simulators have been used for similar research, such as SimGrid [9], GridSim [10], or Alea [11], this experiment needed a simple discrete event simulator using the BFBP scheduler. The simulators mentioned above were either deemed overly complex for our purposes, or they failed to allow for the three limiting resources (memory, CPUs, and

Algorithm 1 Best Fit Bin Packing Scheduling

- 1: while The simulation is incomplete do
- 2: **if** Some job in the queue can be executed on some machine **then**
- 3: Find the (job, machine) pairing which results in the fewest remaining resources for some machine. Begin executing that job on that machine.
- 4: **else**
- 5: Advance simulation time until a new job is submitted or a running job ends, whichever is sooner.
- 6: Queue submitted jobs and stop ending jobs.
- 7: **end if**
- 8: end while

Fig. 1. Pseudocode for the best fit bin packing algorithm

GPUs) we were interested in investigating. An HPC scheduler simulator was also considered, such as the Slurm simulator developed at SUNY University in Buffalo [12]. Although this option was investigated further, scaling a job's actual duration from the log data to the new machine once it is assigned to a machine was challenging. As such, a custom discrete event simulator was developed and utilized for this research. The simulator allows for three resource constraints in each machine: memory, CPUs, and GPUs. It is fairly lightweight, fast, and easy to understand.

A significant consideration when evaluating new server hardware is the performance increase newer technology or architectures can provide. Using log data, we know how long a job took on a machine with known hardware. Since the specifications for the new hardware under consideration are also known, the actual duration of the jobs from the historic log data was scaled using base performance of the processor as reported by SPEC CPU2017 benchmark, second quarter, 2023 [13].

Knowing how a particular job performed on one set of hardware and estimating how it will perform on some other hypothetical set of hardware is challenging. Sharkawi et al. [14] successfully used a similar SPEC benchmark to estimate the performance projections of HPC applications. Other researchers, like Wang et al. [15] have pointed out that these benchmarks fail to account for all the variables affecting job resource utilization and should be avoided. Although CPU performance is not the only factor by which we could have scaled job duration, and perhaps it is not the best factor by which to scale, it worked well for our purposes. The discrete event simulator was implemented such that the scaling factor could be easily changed if other researchers should find a different factor more relevant to their situation.

Various metrics are typically used when evaluating the performance of HPC scheduling algorithms. Some of these are average wait time, HPC utilization, average turnaround time, makespan, throughput, etc. Which metric is used depends on the application and function of the HPC system, and different organizations may value one metric over another. The metric used for this research was average wait time, or the average number of seconds each job spent waiting in the job queue for execution on HPC resources. We presume that the same techniques could be applied by other researchers using a different metric, should they prefer a different one.

Recommender systems power a variety of applications like search engines and music recommendation systems. First, the hits for the system must be defined. Hits are the elements from the data set that are relevant to the user's search. Next, the user specifies the number of recommendations, k, that they would like to receive. If the recommender system is precise, a large portion of the k items returned will be hits.

III. METHODOLOGY

The general plan for optimizing a hardware package for our fixed budget can be summarized as follows:

- 1) Receive vendor quotes with potential server options.
- Generate potential server combinations to purchase under the specified budget which meet our procurement requirements.
- 3) Identify a typical set of jobs representing the workloads typically submitted to our HPC system.
- Conduct simulations using a subset of the server packages to schedule the representative job set and compute metrics to determine their performances.
- 5) Use machine learning to train and refine a model that can predict the performance of un-simulated server combinations.
- 6) Develop a recommender system using the machine learning model.
- 7) Subjectively evaluate the recommended server packages and make a more informed procurement decision.

A. Generate Server Options

To begin, we received several vendor quotes specifying the costs and capabilities of 21 potential servers to purchase. When considering upgrade options, we typically separate servers into one of three categories: compute nodes, big memory nodes, or GPU nodes. A compute node typically has a large number of processor cores, a moderate amount of memory, and no GPU. A big memory node will have a large amount of memory with a moderate amount of CPU cores and no GPU. A GPU node is any node which has a GPU. Table I lays out the options we received from several different vendors. The procurement budget was fixed at \$1 million, and all possible server combinations were generated in the following way:

- Separate servers into three categories: compute nodes, big memory nodes, and GPU nodes.
- Choose all combinations of one node from each category.
- Determine all quantities of the three node types under a given budget such that there is at least one GPU node and there is not enough funding remaining to purchase another node.

TABLE I. SERVER CAPABILITIES AND COSTS UNDER INVESTIGATION

Node type	Distinct	Memory	CPUs	GPUs	Cost
	nodes	range	range	per	range
	consid-	per	per	node	per
	ered	node	node		node
Compute	4	256-512	24-64	0 GPUs	\$6k-
		Gb	cores		\$10k
Big memory	2	1024	24-64	0 GPUs	\$11k-
		Gb	cores		\$13k
GPU	15	256-	24-64	1-8	\$14k-
		1024	cores	GPUs	\$100k
		Gb			

Number of Jobs Submitted over time



Fig. 2. The number of jobs submitted over time for the selected day

In our selected job set, many jobs requested GPUs as a resource. These jobs would automatically fail if at least one GPU node were not included in a potential server package. Roughly 127,000 different server combinations met these requirements. Table II provides an illustrative example of how the server combinations were generated. Many server options and packages were omitted from the table for the sake of brevity.

B. Identify a Representative Set of Jobs

One typical day's worth of submitted jobs (roughly 16,000 jobs) was subjectively pulled from the log data of the local HPC system. As with most HPC systems, jobs were submitted in a bursty manner, and variety of resources were requested. Figure 2 and Table III display some descriptive statistics and information about the jobs used by this research.

C. Job Duration Scaling

The submitted jobs were scaled using the base performance of the processor on the SPEC CPU2017 benchmark suite. The requested duration was not modified, but the actual duration of each job was calculated using the following formula:

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New duration = \frac{\log ged \ duration * \log ged \ processor \ performance}{new \ processor \ performance}
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D. Discrete Event Simulator

A discrete event simulator was implemented in Python that provides the following functionality:

ComputeNode1, \$6,960 ea.	BigMemNode1, \$11,112 ea.	GPUNode1, \$14,730 ea.		Package Cost	Funds Remaining
141	0	1		\$996,090	\$3,910
139	1	1		\$993,282	\$6,718
138	2	1		\$997,434	\$2,566
				•	
:	:	:	· ·	:	:
0	1	67		\$998,022	\$1,978

TABLE III. Descriptive statistics for the pool of selected jobs

	Requested	Requested	Requested	Requested	Actual
	Mem (in	CPUs	GPUs	Duration	Duration
	Gb)			(in	(in
				hours)	hours)
Mean	5.12	4.75	0.002	2.82	2.27
Std Dev.	16.73	3.33	0.055	1.02	13.67
Min	1	1	0	0	0
Max	800	64	4	11.20	11.20

- A global clock to keep track of simulation time.
- Several queues, priority queues, or lists to track jobs as they progress through the execution process: future jobs, queued jobs, running jobs, completed jobs, and unrunnable jobs.
- Jobs and machines are specified using comma separated value (csv) files, which is loaded prior to the simulation.
- Machines have three limiting resources: available memory, CPUs, and GPUs.
- Jobs are specified with the following attributes: submit time, actual duration, and requested duration, memory, CPUs, and GPUs. Jobs track their start time and end time as the simulation progresses to allow for metric calculation.
- Job end time is set when the job starts running as the job start time plus the job actual duration.
- When a job starts running on a machine, that machine's available resources are decremented by the resources requested by the job. Conversely, when a job completes, the machine executing it has its available resources increased by the amount requested by the ending job.
- Jobs with a submit time greater than the current global clock reside in the future jobs priority queue.
- Jobs with a submit time less than or equal to the current global clock, but not yet assigned to a machine, reside in the job queue.
- Jobs that have begun their execution and have an ending time less than the current global clock, reside in the running jobs priority queue.
- Jobs with an ending time less than or equal to the current global clock reside in the completed jobs list.
- If no node in the cluster has adequate resources to run a particular job, that job is moved to the unrunnable jobs list.
- In the event that no queued jobs can run on available resources, the simulation time "fast forwards" to the next event: either job submission or job ending.
- Jobs in the job queue are run as soon as there are available

resources and are chosen using the best fit bin packing scheduling algorithm described in Algorithm 1.

 Actual job duration from logged job data can be scaled to allow for hardware improvement with newer hardware.

E. Machine Learning

Although each simulation completed fairly quickly, requiring no more than 30 minutes each, this particular combination of server quotes yielded roughly 127,000 combinations that need to be evaluated. To reduce the computational requirement, every tenth line from the file with the server combinations was sampled, and roughly 12,700 simulations for these server packages were completed in parallel using HPC resources. By sampling from the generated server packages uniformly, various quantities of each server under consideration were included in the simulated data. Each server package was summarized into the package total memory, total CPUs, and total GPUs, by summing the resources of every machine comprising the package. The average wait time for the simulation served as the label for each package. The data was split into 90% training and 10% test data, and an XGBoost [16] regression model was trained using training data. The regression model was evaluated using root mean squared error (RMSE) on the test data. An accurate regression model enabled the prediction of the average wait time for unsimulated server combinations and saved countless hours of additional simulation.

F. Recommender System

In our case, a hit was defined as a server combination with an average wait time in the lowest 5% of simulated combinations (or 632 hits out of the \sim 12,700 simulated server combinations). The value of k was varied to evaluate the performance of the recommender system. Then, once confidence was gained that our recommender system was functioning properly, it was used to recommend systems from the entire server combination pool of 127,000 server combinations. The recommendations were summarized and evaluated subjectively before arriving at a final procurement decision.

G. Simplifying Assumptions

The current nodes comprising the HPC system were not added to the set of nodes simulating the selected jobs. The benefit current nodes would provide to the new servers under investigation would be common to all.

Any additional equipment required to install and operate the new servers (e.g., networking hardware, additional cooling equipment, server racks, power infrastructure, etc.) were not deducted from the total procurement budget. It was thought that these costs would be a relatively fixed regardless of the server package chosen. The same analysis described in this research could be done by reducing the total budget by the cost of additional hardware and then completing the analysis with a reduced budget.

IV. EVALUATION

Pearson's Correlation Coefficient [17] determined the extent of the correlation between the total memory, CPUs, and GPUs of a package and the average wait time. This coefficient provides a value between -1 and 1, where values closer to -1 or 1 indicate that the feature and the label are more strongly correlated. A coefficient of 0 indicates no correlation.

Wait time was calculated by analyzing the completed jobs output from each simulation. The wait time for each job was the number of seconds from the time the job was submitted until it began. For N jobs, the average wait time was calculated as follows:

AvgWaitTime =
$$\frac{\sum_{i=0}^{N} (\text{Start Time}_i - \text{Submit Time}_i)}{N}$$

Root Mean Squared Error was utilized for regression model evaluation calculated according to the following formula:

$$\text{RMSE} = \sqrt{\frac{\sum_{i=0}^{N} (\text{actual wait time}_i - \text{predicted wait time}_i)^2}{N}}$$

The performance of the final recommender system was evaluated using precision@k, recall@k, and F1@k. In general, precision@k is the proportion of recommended items in the top-k set that are relevant, and recall@k is the proportion of relevant items found in the top-k recommendations. F1@k is the harmonic mean of precision@k and recall@k, which simplifies them into a single metric. They were calculated according to the following formulas:

$$\begin{aligned} & \text{Precision} @k = \frac{(\text{# of recommended items @k that are relevant})}{(\text{# of recommended items @k)}} \\ & \text{Recall} @k = \frac{(\text{# of recommended items @k that are relevant})}{(\text{total # of relevant items})} \\ & \text{F1}@k = \frac{(2*\text{precision}@k*\text{recall}@k)}{(\text{precision}@k+\text{recall}@k)} \end{aligned}$$

V. RESULTS

The correlation of features, the performance of the regression model and the recommender system, and some analysis about the recommended server compositions are described below.

A. Feature Correlation

The correlation between the features and the labels is shown in Table IV. For this set of jobs, the total CPUs in a server package were most strongly correlated to the average wait time. For the chosen jobs, the more CPUs a package had, the lower its average wait time.

Since we are constrained by our available budget of \$1 million, choosing to buy one type of node over another is

a zero-sum game. The more GPU nodes we purchase, and the more GPUs there are per node, the fewer compute nodes or big memory nodes we are able to afford. This is indicated by the positive correlation between GPUs and the average wait time.

TABLE IV. Pearson Correlation Coefficients

	TotalMem	TotalCPUs	TotalGPUs	AvgWaitTime
TotalMem	1.00	0.14	-0.54	-0.23
TotalCPUs	0.14	1.00	-0.42	-0.70
TotalGPUs	-0.54	-0.42	1.00	0.44
AvgWaitTime	-0.23	-0.70	0.44	1.00

B. Regression Model

The XGBoost regression model had a RMSE = 150.13 seconds, indicating that the total memory, CPUs, and GPU features made excellent predictors for the average wait time for these jobs when simulated with the discrete event simulator. The predicted vs. actual wait time is shown in Figure 3. If the regression model were perfect, all these points would lie upon the y = x line, and it is clear that this model does a good job at predicting the average wait time for a given composition of servers.

Predicted vs. Actual Wait Time



Fig. 3. The predicted vs. actual wait times showing the accuracy of our regression model.

C. Recommender System

The regression model was used to predict the 12,700 labeled simulations, and their precision@k, recall@k, and F1@k for various values of k are displayed in Table V. The goal was to reduce the number of possibilities from roughly 127,000 different possible combinations of servers down to a reasonable number which could be evaluated by an HPC system administrator and have a large percentage of the recommended server combinations be hits (among the best 5% of server combinations with the lowest average wait times). Although precision@10 was 100%, it is thought that seeing more server package options would allow system administrator could easily and quickly review up to 50 recommendations (k = 50),

and more than 46 out of 50 of these recommendations returned by this system (92%) would be top performing server combinations, which is excellent. Recall@k when k is less than the number of total hits (632 hits total) is unfairly penalized, but the recall@k above 632 is also excellent. When k = 1,000, the recall@1000 = 91%, meaning the recommender system successfully retrieved 91% of the top 5% performing server packages when returning less than 1% of the 127,000 different options.

TABLE V. Precision@k and Recall@k for Test Data

k value	Precision@k	Recall@k	F1@k
10	1.00	0.02	0.03
50	0.92	0.07	0.13
100	0.81	0.13	0.22
500	0.74	0.59	0.66
632	0.72	0.72	0.72
1000	0.58	0.91	0.71

D. Recommended Compositions

Beyond looking at the individual server compositions recommended, we wanted to draw some conclusion about the types and quantity of nodes that the recommender system returned. The sum of the server quantities for the top 50 recommendations can be found in Table VI. Compute nodes with the larger number of cores were vastly preferred, and the recommender system did not recommend spending additional funds on more memory for the compute nodes. Additionally, the recommender system preferred the cheaper big memory node with fewer cores. Finally, for our typical workload, the recommender system did not recommended purchasing a large number of GPUs per GPU node, instead recommending servers with 2 GPUs per server most often. In terms of our budget breakdown, the recommender system suggests spending on average 58% of our total budget on compute nodes, 8% on big memory nodes, and 34% on GPU nodes.

TABLE VI. Recommendations Drawn from Model Predicted Results

Node Type	Node Description	Sum of Servers Across
Noue Type	Node Description	0
		<i>Top 50</i>
	Cheapest w/ 256Gb	232
Compute Nodes	Cheapest w/ 512Gb	0
	Expensive w/ 256Gb	3,467
	Expensive w/ 512Gb	0
Big Memory Nodes	Cheapest w/ 1024Gb	232
Big Memory Nodes	Expensive w/ 1024Gb	111
GPU Nodes	2 GPUs in one server	732
	4 GPUs in one server	267

VI. CONCLUSIONS

This recommender system is not intended to replace the expertise of HPC administrators when it comes to decisions for hardware procurement. It is our hope that this tool can provide a data-driven technique which will help narrow the search space with which administrators are confronted when they make procurement decisions. Returning to the research question: experimental simulation coupled with a regression model enabled a recommender system to return server compositions under a given budget with low average wait times with a precision@50 of 92%. Additionally, the discrete event simulator, job data set, machine learning code, and recommender system code are released under the GPLv3 license should other researchers find it useful (https://github.com/shutchison/ Optimal-Hardware-Procurement-for-a-HPC-Expansion).

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